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Natural extension of choice functions

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Goal of the paper



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- Coherent choice functions can be used as a model of the rational behaviour of an individual or a group.
- They were extended by Seidenfeld et al. to allow for incomparability, that arises naturally with imprecise information.
- Previous works assume that the choice function is determined for all options, something unreasonable in practice.
- Given a partially specified choice function, can we determine its implications on other option sets, using *only* the axioms of coherence?

Choice functions



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We consider a real vector space V that represents our option space, and let Q be the set of all non-empty *finite* subsets of V.

A choice function C on $\mathcal V$ is a map

$$C \colon \mathcal{Q} \to \mathcal{Q} \cup \{\emptyset\} \colon A \mapsto C(A) \text{ such that } C(A) \subseteq A.$$

Equivalently to a choice function C, we may consider its rejection function R, defined by $R(A) := A \setminus C(A)$ for all A in \mathcal{Q} .

We will assume that \mathcal{V} is ordered by a vector ordering \leq , and that \prec is its associated strict partial order \prec .

Coherent choice functions



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We call a rejection function R on \mathcal{V} coherent if for all A, A_1 and A_2 in \mathcal{Q} , all u and v in \mathcal{V} , and all λ in $\mathbb{R}_{>0}$:

- R1. $R(A) \neq A$;
- R2. if $u \prec v$ then $u \in R(\{u, v\})$;
- R3. a. if $A_1 \subseteq R(A_2)$ and $A_2 \subseteq A$ then $A_1 \subseteq R(A)$; b. if $A_1 \subseteq R(A_2)$ and $A \subseteq A_1$ then $A_1 \setminus A \subseteq R(A_2 \setminus A)$;
- R4. a. if $A_1\subseteq R(A_2)$ then $\lambda A_1\subseteq R(\lambda A_2)$; b. if $A_1\subseteq R(A_2)$ then $A_1+\{u\}\subseteq R(A_2+\{u\})$.



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Natural extension: definition



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Let \mathcal{Q}_0 denote those option sets that include 0. Given an assessment $\mathcal{B} \subseteq \mathcal{Q}_0$, it has the interpretation that 0 should be rejected from every option set B in \mathcal{B} .

The natural extension of \mathcal{B} is the rejection function

$$\mathcal{E}(\mathcal{B}) \coloneqq \inf\{R \text{ coherent} : (\forall B \in \mathcal{B})0 \in R(B)\}$$
$$= \inf\{R \text{ coherent} : R \text{ extends } \mathcal{B}\},$$

where we let $\inf \emptyset$ be equal to $id_{\mathcal{Q}}$, the identity rejection function that maps every option set to itself.

An operational definition



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For any set of options A in \mathcal{Q} , let $R_{\mathcal{B}}(A)$ be given by

$$\Big\{u \in A: (\exists A' \in \mathcal{Q}) \Big(A' \supseteq A \text{ and } (\forall v \in \{u\} \cup (A' \setminus A))$$
$$\Big((A' - \{v\}) \cap \mathcal{V}_{\succ 0} \neq \emptyset \text{ or } (\exists B \in \mathcal{B}, \exists \mu \in \mathbb{R}_{> 0}) \{v\} + \mu B \preccurlyeq A'\Big)\Big)\Big\},$$

where $\mathcal{V}_{\succ 0} \coloneqq \{u \in \mathcal{V} : 0 \prec u\}.$

- $R_{\mathcal{B}}$ is the least informative rejection function that satisfies Axioms R2–R4 and extends \mathcal{B} .
- We say that $\mathcal B$ avoids complete rejection when $R_{\mathcal B}$ satisfies Axiom R1.

Characterizing the natural extension



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For any assessment $\mathcal{B} \subseteq \mathcal{Q}_0$, the following are equivalent:

- (i) \mathcal{B} avoids complete rejection;
- (ii) ${\cal B}$ has a coherent extension;
- (iii) $\mathcal{E}(\mathcal{B}) \neq \mathrm{id}_{\mathcal{Q}}$;
- (iv) $\mathcal{E}(\mathcal{B})$ is coherent;
- (v) $\mathcal{E}(\mathcal{B})$ is the least informative rejection function that is coherent and extends \mathcal{B} .

When any of these equivalent statements hold, then $\mathcal{E}(\mathcal{B}) = R_{\mathcal{B}}$.



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Binary comparisons: sets of desirable options



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A desirability assessment $B \subseteq \mathcal{V}$ is a set of options that we strictly prefer to the zero option.

We call a set of desirable options $D \subseteq \mathcal{V}$ coherent if for all u and v in \mathcal{V} and λ in $\mathbb{R}_{>0}$:

D1. $0 \notin D$;

D2. if $0 \prec u$ then $u \in D$:

D3. if $u \in D$ then $\lambda u \in D$:

D4. if $u, v \in D$ then $u + v \in D$.

Desirable options and choice functions



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There is a one-to-one correspondence between sets of options $B \subseteq \mathcal{V}$ and sets of binary comparisons: $\mathcal{B}_B := \{\{0, u\} : u \in B\}.$

With this correspondence, if D is a coherent set of desirable options D, then the rejection function R_D given by

$$R_D(A) = \{ u \in A : (\forall v \in A)v - u \notin D \}$$

for all A in \mathcal{Q} is coherent.

More generally, given $B \subseteq \mathcal{V}$, we say that $D \subseteq \mathcal{V}$ extends B if $B \subseteq D$.

D extends $B \Leftrightarrow R_D$ extends \mathcal{B}_B .

Natural extension of sets of desirable options



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Consider any desirability assessment $B \subseteq \mathcal{V}$. We say that it avoids non-positivity when it is included in some coherent set of desirable options. In that case, the smallest such set is its natural extension

$$\mathcal{E}^{\mathbf{D}}(B) := \inf\{D \text{ coherent } : B \subseteq D\} = \operatorname{posi}(\mathcal{V}_{\succ 0} \cup B),$$

where we let $\inf \emptyset = \mathcal{V}$.

- B avoids non-positivity $\Leftrightarrow \mathcal{B}_B$ avoids complete rejection.
- In that case, $\mathcal{E}(\mathcal{B}_B) = R_{\mathcal{E}^{\mathbf{D}}(B)}$.

Connection between the natural extensions

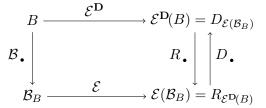


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Implication: not all coherent choice functions are determined by binary comparisons

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Many examples of coherent choice functions, such as the M-admissible or E-admissible ones can be written as the infima of choice functions determined by binary comparisons.

By considering $B := \{0, f, \lambda f\}$ with f a gamble and λ an element of $\mathbb{R}_{>0}$ and different from 1, we obtain that its natural extension $R_{\mathcal{B}}$ is coherent but is NOT the infima of binary choice functions.

A consequence of this is that either (i) choice functions do NOT form a strong belief structure, or (ii) there are maximal (=maximally informative) choice functions that are NOT determined by binary comparisons.



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Sets of indifferent options



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In addition to the set D of options that we prefer to the status quo, we can also consider the set I that we consider indifferent to it. We say that this set is coherent if for all u,v in $\mathcal V$ and λ in $\mathbb R$:

 $I_1. \ 0 \in I;$

 I_2 . if $u \in \mathcal{V}_{\succ 0} \cup \mathcal{V}_{\prec 0}$ then $u \notin I$;

 I_3 . if $u \in I$ then $\lambda u \in I$;

 I_4 . if $u, v \in I$ then $u + v \in I$.

Compatibility with choice functions



A coherent set of indifferent options I determines a quotient space

 $V/I := \{[u] : u \in V\} = \{\{u\} + I : u \in V\} = \{u/I : u \in V\}.$

R is compatible with I if there exists R' on $\mathcal{Q}(\mathcal{V}/I)$ such that

$$R(A) = \{ u \in A : [u] \in R'(A/I) \} \ \forall A \in \mathcal{Q}(\mathcal{V}).$$

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Compatibility with natural extension



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Given $\mathcal{B} \subseteq \mathcal{Q}_0(\mathcal{V})$ and any coherent set of indifferent options I, the natural extension of \mathcal{B} under I is the rejection function

$$R_{\mathcal{B},I}(A) \coloneqq \{u \in A : [u] \in R_{\mathcal{B}/I}(A/I)\} \text{ for all } A \text{ in } \mathcal{Q}(\mathcal{V}),$$

where
$$\mathcal{B}/I \coloneqq \{B/I : B \in \mathcal{B}\} \subseteq \mathcal{Q}_{[0]}(\mathcal{V}/I)$$
.

 $R_{\mathcal{B},I}$ is the least informative rejection function that is coherent, extends \mathcal{B} , and is compatible with I, if there is one.



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Conclusions and open problems



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Conclusions:

- The notion of natural extension can be extended to the theory of coherent choice functions.
- Binary comparisons (=sets of desirable options) follow as a particular case.
- Coherent choice functions are not a strong belief structure.

Open problems:

References



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